Safety in multi-assembly via paths appearing in all path covers of a DAG

Manuel Cáceres, Brendan Mumey, Edin Husić, Romeo Rizzi, Massimo Cairo, Kristoffer Sahlin, Alexandru I. Tomescu

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Many problems in bioinformatics can be seen as puzzles



Genome Assembly

Reconstruct a genomic sequence based on reads obtained from it.



Multi-Assembly

Reconstruct multiple sequences based on *mixed-reads* obtained from all of them.



Puzzles solved perfectly



Unsolvable puzzles (multiple solutions)



Puzzles with multiple solutions



Safe parts common to *all* possible solutions



Contigs in Genome Assembly

Genomic fragments *promised* to occur in the original genome.

▶ Used and reported by practical assemblers [10, 11, 17, 20].

Developed theoretically [24, 6, 7, 5].

Completeness

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Completeness

Not used in multi-assembly!

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- Directed acyclic graph.
- E.g.:
 RNA transcript assembly [25, 21, 3, 13, 9, 16, 19, 4, 12, 23, 22, 14].
 Viral quasi-species assembly [8, 27, 2, 1]

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Minimum path cover (MPC).

- Of size *k*, known as the width.
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Generalized path cover.

- Of size $\leq \ell$ (parameter).
- Paths starting at $S \subseteq V$ (parameter).
- Paths ending at $T \subseteq V$ (parameter).

Our results

Theoretical results

Theoretical results

	Minimum path covers	Generalized path covers
Safe edges	O(k V E)	O(k V E)
Maximal safe paths	$O(k^2 V E)$	$O(max(1, 2k - \ell)k V E)$

Recall k is the width of the graph (size of an MPC)

Practical results

Apply to RNA transcript assembly defining RNA contigs.

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Apply to RNA transcript assembly defining RNA contigs.

Proof-of-concept study

- Splicing graphs from human annotated transcripts.
- Double length compared to unitigs.
- Transcript coverage of 80%.
- Less than 15 seconds (all transcript annotation from Ensembl database [26]).

Theoretical results (MPC only)

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- 2. For each subpart \rightarrow Every subpath *P* of a path of \mathcal{P} .

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▶ In O(|V||E|) [18].

2. For each subpart \rightarrow Every subpath *P* of a path of \mathcal{P} .

• Test if P is safe \rightarrow Compute the width of G^{P} .

Safe edges

Theorem (Safe edges MPC) An edge e is safe if and only if width(G) < width($G \setminus e$).

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- Naive algorithm: $O(k|V|^2|E|)$.
- But, we can do better.

Shrinking Primitive

Lemma (Shrinking [15])

Given a path cover $\mathcal{P} = P_1, \ldots, P_t$ of G, we can obtain a MPC of G in time O(|E|(t - k + 1))

Theorem (Safe edges MPC)

An edge e is safe if and only if width(G) < width(G $\setminus e$).

• Naive algorithm: $O(k|V|^2|E|)$.

But, we can do better.

$$\blacktriangleright O(\sum_{e \in \mathcal{P}} \mu_e |E|) = O(k|V||E|)$$

Maximal safe paths

Path Reduction

Definition

Given a path $P = x_1, \ldots, x_p$ of G, we define $G^P = (V, E^P)$, where

$$E^{P}=(E\setminus\{(x_{p-1},x_{p})\})\cup\bigcup_{i=2}^{p}\left\{(u,x_{p})\mid u\in N^{-}(x_{i})\setminus\{x_{i-1}\}\right\}.$$

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Theorem (Maximal safe paths MPC)

Let $P = x_1, ..., x_p$ be a path of G, such that $x_1, ..., x_{p-1}$ is a safe path. It holds that P is safe if and only if width $(G) < \text{width}(G^P)$.

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• Paths of increasing size algorithm: $O(k|V|^3|E|)$.

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- Paths of increasing size algorithm: $O(k|V|^3|E|)$.
- By using shrinking, and two-finger on each path: $O(k^2|V||E|)$.

Practical results (RNA contigs)

Central dogma of molecular biology















RNA contigs



RNA contigs = Maximal safe paths for generalized path cover

Experimental setup

Human gene annotation (Ensembl [26],GRCh38.p13).

All transcripts in chromosomes 1 to 22, on the forward strand.

Build splicing graph.

- Perfect scenario.
- Known answers.

$$\blacktriangleright \ \ell \in \{k, k+1, \ldots, 2k\}.$$

▶ Baseline comparison with *ST*-unitigs.

Results

(less than 15 seconds to run)

Results

	small graphs (3-21 vertices)		medium graphs (22-52 vertices)		large graphs (53-725 vertices)	
l	prec	mcr	prec	mcr	prec	mcr
k	0.84	0.82	0.81	0.64	0.84	0.56
	1.00	0.86	0.92	0.62	0.89	0.53
k+1	1.00	0.79	0.99	0.61	0.99	0.53
	1.00	0.83	1.00	0.59	1.00	0.50
t	1.00	0.79	1.00	0.61	1.00	0.53
	1.00	0.83	1.00	0.59	1.00	0.50
2 <i>k</i>	1.00	0.79	1.00	0.61	1.00	0.53
	1.00	0.83	1.00	0.59	1.00	0.50
ST-	1.00	0.64	1.00	0.49	1.00	0.42
unitigs	1.00	0.67	1.00	0.47	1.00	0.39

Precision and Relative Maximum Coverage (of transcripts) for RNA contigs and *ST*-unitigs.

Results

l	small graphs (3-21 vertices)	medium graphs (22-52 vertices)	large graphs (53-725 vertices)
k	2.65×	3.54×	3.50×
k+1	1.43×	1.83×	1.97×
t	1.42×	1.82×	1.94×
2 <i>k</i>	1.42×	1.82×	1.94×

Relative length of longest RNA contig containing a ST-unitig.

 Efficient algorithms obtaining all maximal safe paths for generalized path covers.

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 - Publicly available code, datasets, and manipulation scripts.
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- Proof-of-concept RNA contigs on human annotated transcripts.
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- Future developments of safe paths
 - Output of transcript assemblers.
 - Use as preprocessing step.
 - Validate transcript assemblies.
 - Apply to other multi-assembly problem.

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